

2.098/6.255/15.093 - Recitation 10

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1 Karush Kuhn Tucker Necessary Conditions

$$\begin{aligned} \text{P: } \min \quad & f(x) \\ \text{s.t. } \quad & g_j(x) \leq 0, \quad j = 1, \dots, p \\ & h_i(x) = 0, \quad i = 1, \dots, m \end{aligned}$$

Theorem. (KKT Necessary Conditions for Optimality)

If \hat{x} is local minimum of P and the following *Constraint Qualification Condition (CQC)* holds:

- The vectors $\nabla g_j(\hat{x})$, $j \in \mathcal{I}(\hat{x})$ and $\nabla h_i(\hat{x})$, $i = 1, \dots, m$, are linearly independent, where $\mathcal{I}(\hat{x}) = \{j : g_j(\hat{x}) = 0\}$ is the set of indices corresponding to active constraints at \hat{x} .

Then, there exist vectors (u, v) s.t.:

1. $\nabla f(\hat{x}) + \sum_{j=1}^p u_j \nabla g_j(\hat{x}) + \sum_{i=1}^m v_i \nabla h_i(\hat{x}) = 0$
2. $u_j \geq 0$, $j = 1, \dots, p$
3. $u_j g_j(\hat{x}) = 0$, $j = 1, \dots, p$ (or equivalently $g_j(\hat{x}) < 0 \Rightarrow u_j = 0$, $j = 1, \dots, p$)

Theorem. (KKT + Slater)

If \hat{x} is local minimum of P and the following *Slater Condition* holds:

- There exists some feasible solution \bar{x} such that $g_j(\bar{x}) < 0$, $\forall j \in \mathcal{I}(\hat{x})$, where $\mathcal{I}(\hat{x}) = \{j : g_j(\hat{x}) = 0\}$ is the set of indices corresponding to active constraints at \hat{x} .

Then, there exist vectors (u, v) s.t.:

1. $\nabla f(\hat{x}) + \sum_{j=1}^p u_j \nabla g_j(\hat{x}) + \sum_{i=1}^m v_i \nabla h_i(\hat{x}) = 0$
2. $u_j \geq 0$, $j = 1, \dots, p$

3. $u_j g_j(\hat{x}) = 0, j = 1, \dots, p$ (or equivalently $g_j(\hat{x}) < 0 \Rightarrow u_j = 0, j = 1, \dots, p$)

Example.

Solve

$$\begin{aligned} \min \quad & x_1^2 + x_2^2 + x_3^2 \\ \text{s.t.} \quad & x_1 + x_2 + x_3 \leq -18 \end{aligned}$$

Solution.

For some x to be a local minimum, condition (1) requires that $\exists u$ s.t. $2x_i + u = 0, i = 1, 2, 3$.

Now, the constraint can either be active or inactive:

- If it is inactive, then $u < 0$ by condition (3). This would imply $x_1 = x_2 = x_3 = 0$, but $x = (0, 0, 0)^\top$ is infeasible, so this cannot be true of a local minimum.
- If it is active, then $x_1 + x_2 + x_3 = -18$ and $2x_i + u = 0, i = 1, 2, 3$. This is a system of four linear equations in four unknowns. We solve and obtain $u = 12, x = (-6, -6, -6)^\top$. Since $u = 12 \geq 0$, there exists a u as desired. To check the regularity requirement, we simply confirm that $\nabla x = (1, 1, 1)^\top \neq 0$. Also, we could have checked that the Slater condition is satisfied (eg use $\bar{x} = (-10, -10, -10)^\top$).

Hence $(-6, -6, -6)^\top$ is the only candidate for a local minimum. Now, the question is: is it a local minimum? (Note that since this is the unique candidate, this is the same as asking if the function has a local minimum over the set.)

Observe that the objective function is convex. Why? Because it is a positive combination of convex functions. Now, is the feasible set convex? Answer: yes, since it is of the form $\{x \in \mathbb{R}^n : f(x) \leq 0\}$, where f is a convex function.

So we may apply a stronger version of the KKT conditions, the KKT *sufficient* conditions, which imply that any x which satisfies the KKT necessary conditions and also meets these two convexity requirements is in fact a global minimum.

So the point $x = (-6, -6, -6)^\top$ is the unique global minimum (unique since it was the only candidate).

Example.¹

Solve

$$\begin{aligned} \min \quad & -\log(x_1 + 1) - x_2 \\ \text{s.t.} \quad & g(x) \triangleq 2x_1 + x_2 - 3 \leq 0 \\ & x_1, x_2 \geq 0 \end{aligned}$$

¹Thanks to Andy Sun.

Solution.

Firstly, observe that this is a convex optimization problem, since $f(x)$ is convex (a positive combination of the convex functions $-x_2$ and $-\log(x_1 + 1)$), and the constraint functions $g(x)$, $-x_1$ and $-x_2$ are convex (again, this is required for the feasible space to be convex, since we have \leq constraints). Alternatively, in this case we can see that the feasible space is a polyhedron, which we know to be convex.

Now, in order to use KKT, we need to assume which inequalities are active. Let's start by assuming that at a local minimum x , only $g(x) \leq 0$ is active. This leads to the system: $\begin{bmatrix} -1 \\ x_1+1 \\ -1 \end{bmatrix} + u \begin{bmatrix} 2 \\ 1 \end{bmatrix} = 0$, which gives $u = 1$ and $x_1 = -0.5$, which is not feasible, so our assumption cannot be correct.

Now try assuming $g(x) \leq 0$ and $-x_1 \leq 0$ are active, giving the system $\begin{bmatrix} -1 \\ x_1+1 \\ -1 \end{bmatrix} + u_1 \begin{bmatrix} 2 \\ 1 \end{bmatrix} + u_2 \begin{bmatrix} -1 \\ 0 \end{bmatrix} = 0$, which gives $u_1 = 1, u_2 = 1$ and $x_2 = 3$ (recall we assumed $x_1 = 0$).

Now since it's a convex optimization problem, and the Slater condition is trivially satisfied, this is a global minimum by the KKT sufficient conditions.

Example.

The following example shows that the KKT theorem may not hold if the regularity condition is violated: Consider

$$\begin{aligned} \min \quad & x_1 + x_2 \\ \text{s.t.} \quad & (x_1 - 1)^2 + x_2^2 - 1 = 0 \\ & (x_1 - 3)^2 + x_2^2 - 9 = 0 \end{aligned}$$

The feasible region is the intersection of two circles, one centered at the point $(1, 0)$ with radius 1, the other at the point $(3, 0)$ with radius 3. The intersection occurs at the origin, which is the optimal solution by inspection.

We have $\nabla f(\hat{x}) = (1, 1)^\top$, $\nabla h_1(\hat{x}) = (2x_1 - 2, 2x_2)^\top = (-2, 0)^\top$, and $\nabla h_2(\hat{x}) = (2x_1 - 6, 2x_2)^\top = (-6, 0)^\top$. So condition (1) cannot hold.

2 Conjugate Gradient Method²

Consider minimizing the quadratic function $f(x) = \frac{1}{2}x^\top Qx + c^\top x$.

1. d_1, d_2, \dots, d_m are **Q-conjugate** if

$$d_i^\top Q d_j = 0, \quad \forall i \neq j$$

²Thanks to Allison Chang for notes

2. Let x_0 be our initial point.
3. Direction $d_1 = -\nabla f(x_0)$.
4. Direction $d_{k+1} = -\nabla f(x_{k+1}) + \lambda_k d_k$, where $\lambda_k = \frac{\nabla f(x_{k+1})^\top d_k}{d_k^\top Q d_k}$ in the quadratic case (and $\lambda_k = \frac{\|\nabla f(x_{k+1})\|^2}{\|\nabla f(x_k)\|^2}$ in the general case). It turns out that with each d_k constructed in this way, d_1, d_2, \dots, d_k are \mathbf{Q} -conjugate.
5. By Expanding Subspace Theorem, x_{k+1} minimizes $f(x)$ over the affine subspace $S = x_0 + \text{span}\{d_1, d_2, \dots, d_k\}$.
6. Hence finite convergence (n steps).

3 Barrier Methods

A barrier function $G(x)$, is a continuous function with the property that it approaches ∞ as one of the $g_j(x)$ approaches 0 from below.

Examples:

$$-\sum_{j=1}^p \log[-g_j(x)] \quad \text{and} \quad -\sum_{j=1}^p \frac{1}{g_j(x)}$$

Consider the primal/dual pair of linear optimization problems

$$\begin{array}{ll} \text{P:} & \min \quad c^\top x \\ & \text{s.t.} \quad Ax = b \\ & \text{s.t.} \quad x \geq 0 \end{array} \qquad \begin{array}{ll} \text{D:} & \max \quad b^\top p \\ & \text{s.t.} \quad A^\top p + s = c \\ & \text{s.t.} \quad s \geq 0 \end{array}$$

To solve P, we define the following barrier problem:

$$\text{BP:} \quad \min \quad B_\mu(x) \triangleq c^\top x - \mu \sum_{j=1}^n \log x_j \\ \text{s.t.} \quad Ax = b$$

Assume that for all $\mu > 0$, BP has an optimal solution $x(\mu)$. This optimum will be unique. Why?

As μ varies, the $x(\mu)$ form what is called the *central path*.

Theorem. $\lim_{\mu \rightarrow 0} x(\mu)$ exists and $x^* = \lim_{\mu \rightarrow 0} x(\mu)$ is an optimal solution to P.

Then the barrier problem from the dual problem is

$$\text{BD:} \quad \max \quad b^\top p + \mu \sum_{j=1}^n \log s_j \\ \text{s.t.} \quad A^\top p + s = c$$

Theorem. Let $\mu > 0$. Then $x(\mu), s(\mu), p(\mu)$ are optimal solutions to BP and BD if and only if the following hold:

$$\begin{aligned} Ax(\mu) &= b \\ x(\mu) &\geq 0 \\ A^\top p(\mu) + s(\mu) &= c \\ s(\mu) &\geq 0 \\ x_j(\mu)s_j(\mu) &= \mu, \forall j \end{aligned}$$

To solve BP using the *Primal path following algorithm*, we:

1. Start with a feasible interior point solution $x_0 > 0$
2. Step in the Newton direction $d(\mu) = (I - X^2 A^\top (A X^2 A^\top)^{-1} A) (X e - \frac{1}{\mu} X^2 c)$
3. Decrement μ
4. Iterate until convergence is obtained (complementary slackness above is ϵ -satisfied)

Note if we were to fix μ and carry out several Newton steps, then x would converge to $x(\mu)$. By taking a single step in the *Newton direction* we can guarantee that x stays “close to” $x(\mu)$, i.e. the *central path*. Hence following the iterative Primal path following algorithm we will converge to an optimal solution by this result and the first theorem above.

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