

## Hints for Problem Set 3

- *Problem 1e*  
Given a matrix  $M$ , how can you make a new matrix  $N$  such that the largest eigenvalue of  $N$  is related to the smallest eigenvalue of  $M$ ?
- *Problem 1f*  
Try to generalize the idea from 1e. Given a matrix  $M$ , how can you create a new matrix  $N$  so that the biggest eigenvalue of  $N$  is related to the eigenvalue of  $M$  closest to  $\lambda$ ?
- *Problem 2b*  
If I have a polynomial  $p(x)$  in one variable and I apply it to a matrix with the given eigenvalue bounds, I get a new matrix  $N$ . What can I say about its eigenvalues? What properties of  $p$  will guarantee that  $N$  still has the same eigenvector with eigenvalue 1 but that every vector orthogonal to it gets shrunk significantly in norm?
- *Problem 2c*  
Look really carefully at the description of conjugate gradient in Shewchuk's article. I claim that we've already computed everything I'm asking you for.
- *Problem 2e*  
For any polynomial  $p$  of degree  $t-1$ , show that there is a vector  $v$  in  $K(A, x, t)$  such that  $v = p(A)x$ . Argue that for a random  $x$ , there is some polynomial  $p$  for which the corresponding  $v$  gives a good Rayleigh quotient, and thus a good estimate for the biggest and smallest eigenvalues of  $A$ .
- *Problem 3*  
Make a new  $2n \times 2n$  matrix  $A$  and a vector  $q$ , and solve the linear system:  
 $Ax = q$   
  
 $q$  will be the vector  $[b; -b]$ .  
 $A$  will be made (somehow) by combining (somehow—maybe you'll need to add them to each other, etc.) the following pieces:  
  
 $D =$  the diagonal of  $M$   
 $P =$  the matrix of positive entries of  $M$  (with everything else set to zero)  
 $N =$  the matrix of negative entries of  $M$  (with everything else set to zero)
- *Problem 4d*  
You have to show how to route the edges of  $G$  over those of  $H'$  with low maximum total stretch, so that you can apply part (a). Part (c) tells you how to route the edges that are internal to a given  $G_i$ . Use this to get a bound on the contribution of these edges to the total stretch.  
  
You then just need to figure out how to route the edges that cross from one  $G_i$  to some other  $G_j$ . Route these edges internally in  $G_i$ , across the bridge edge we've added, and then internally in  $G_j$ . Bound the contribution of these edges.

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