

TENSOR DECOMPOSITIONS AND THEIR APPLICATIONS

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Algorithmic Aspects of Machine Learning

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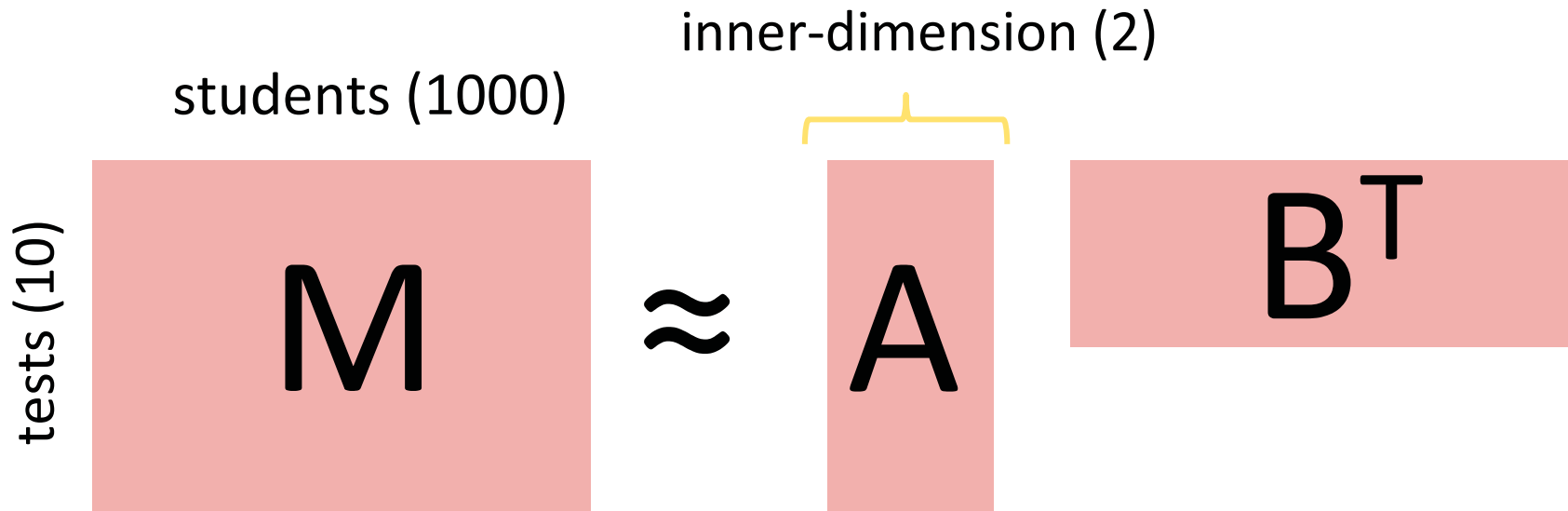
Note: These are unpolished, incomplete course notes.

Developed for educational use at MIT and for publication through MIT OpenCourseware.

SPEARMAN'S HYPOTHESIS

Charles Spearman (1904): There are two types of intelligence, *eductive* and *reproductive*

To test this theory, he invented **Factor Analysis:**



eductive (adj): the ability to make sense out of complexity

reproductive (adj): the ability to store and reproduce information

Given: $M = \sum a_i \otimes b_i$

$$= \underbrace{A B^T}_{\text{"correct" factors}} = \underbrace{A R R^{-1} B^T}_{\text{alternative factorization}}$$

When can we recover the factors a_i and b_i uniquely?

Claim: The factors $\{a_i\}$ and $\{b_i\}$ are not determined uniquely unless we impose additional conditions on them

e.g. if $\{a_i\}$ and $\{b_i\}$ are orthogonal, or $\text{rank}(M)=1$

This is called the **rotation problem**, and is a major issue in factor analysis and motivates the study of **tensor methods**...

OUTLINE

The focus of this tutorial is on Algorithms/Applications/Models for tensor decompositions

Part I: Algorithms

- The Rotation Problem
- Jennrich's Algorithm

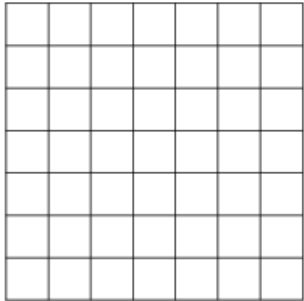
Part II: Applications

- Phylogenetic Reconstruction
- Pure Topic Models

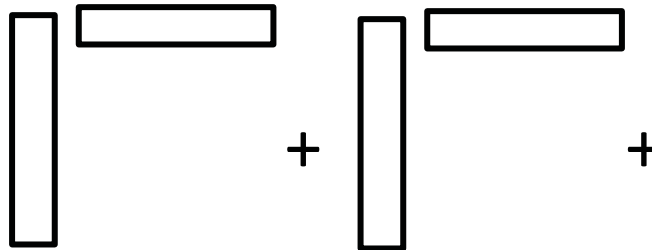
Part III: Smoothed Analysis

- Overcomplete Problems
- Kruskal Rank and the Khatri-Rao Product

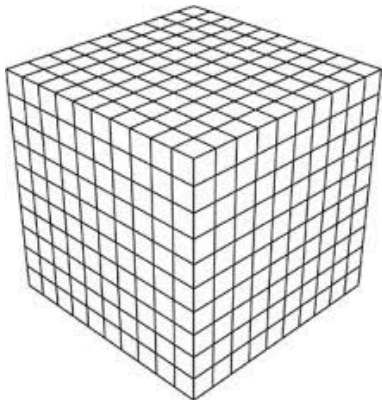
MATRIX DECOMPOSITIONS



$$M = a_1 \otimes b_1 + a_2 \otimes b_2 + \cdots + a_R \otimes b_R$$



TENSOR DECOMPOSITIONS



$$T = a_1 \otimes b_1 \otimes c_1 + \cdots + a_R \otimes b_R \otimes c_R$$

(i, j, k) entry of $x \otimes y \otimes z$ is $x(i) \times y(j) \times z(k)$

When are tensor decompositions unique?

Theorem [Jennrich 1970]: Suppose $\{a_i\}$ and $\{b_i\}$ are linearly independent and no pair of vectors in $\{c_i\}$ is a scalar multiple of each other. Then

$$T = a_1 \otimes b_1 \otimes c_1 + \cdots + a_R \otimes b_R \otimes c_R$$

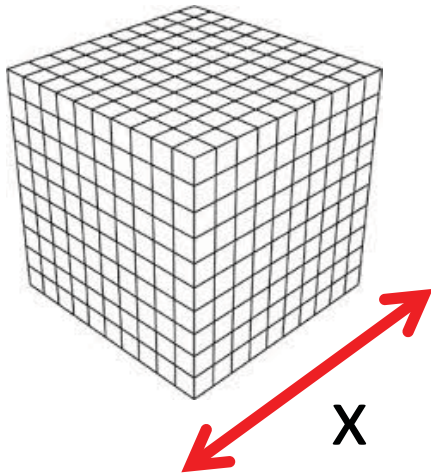
is unique up to permuting the rank one terms and rescaling the factors.

Equivalently, the rank one factors are **unique**

There is a simple algorithm to compute the factors too!

JENNRICH'S ALGORITHM

➔ Compute $T(\bullet, \bullet, x)$

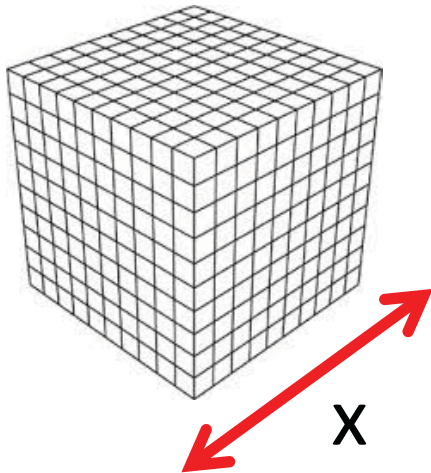


i.e. add up matrix slices

$$\sum x_i T_i$$

JENNRICH'S ALGORITHM

➔ Compute $T(\bullet, \bullet, x)$



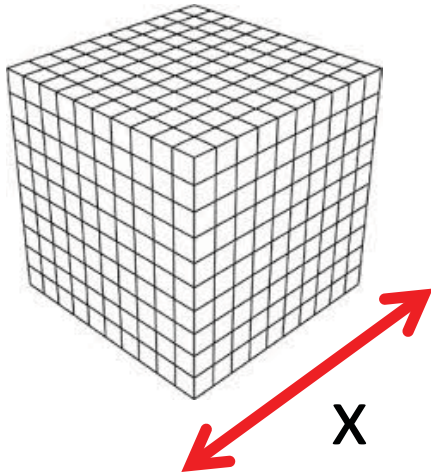
i.e. add up matrix slices

$$\sum x_i T_i$$

If $T = a \otimes b \otimes c$ then $T(\bullet, \bullet, x) = \langle c, x \rangle a \otimes b$

JENNRICH'S ALGORITHM

➔ Compute $T(\bullet, \bullet, x)$

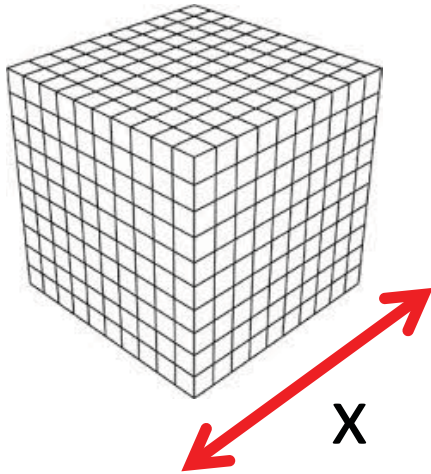


i.e. add up matrix slices

$$\sum x_i T_i$$

JENNRICH'S ALGORITHM

➔ Compute $T(\bullet, \bullet, x) = \sum \langle c_i, x \rangle a_i \otimes b_i$

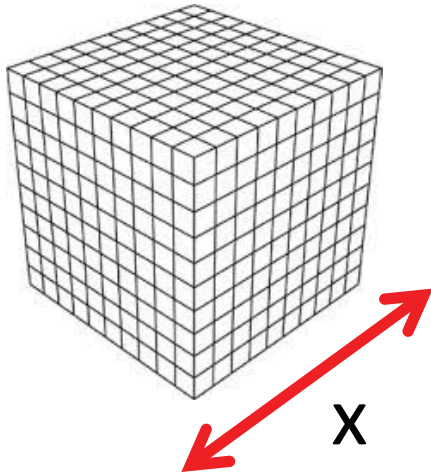


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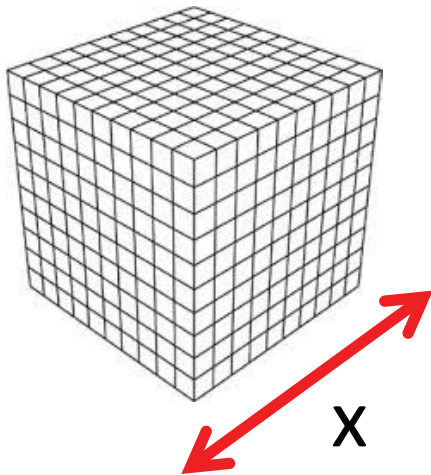
$$\sum x_i T_i$$

(x is chosen uniformly at random from S^{n-1})

JENNRICH'S ALGORITHM

$$\text{Diag}(\langle c_i, x \rangle)$$

➔ Compute $T(\cdot, \cdot, x) = A D_x B^T$



i.e. add up matrix slices

$$\sum x_i T_i$$

(x is chosen uniformly at random from S^{n-1})

JENNRICH'S ALGORITHM

➔ Compute $T(\bullet, \bullet, x) = A D_x B^T$

➔ Compute $T(\bullet, \bullet, y) = A D_y B^T$


➔ Diagonalize $T(\bullet, \bullet, x) T(\bullet, \bullet, y)^{-1}$

JENNRICH'S ALGORITHM

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➔ Diagonalize $T(\bullet, \bullet, x) T(\bullet, \bullet, y)^{-1}$


$$A D_x B^T (B^T)^{-1} D_y^{-1} A^{-1}$$

JENNRICH'S ALGORITHM

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$$A D_x D_y^{-1} A^{-1}$$

Claim: whp (over x, y) the eigenvalues are distinct, so the Eigendecomposition is unique and recovers a_i 's

JENNRICH'S ALGORITHM

➔ Compute $T(\bullet, \bullet, x) = A D_x B^T$

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➔ Match up the factors (their eigenvalues are reciprocals) and find $\{c_i\}$ by solving a linear syst.

Given: $M = \sum a_i \otimes b_i$

When can we recover the factors a_i and b_i uniquely?

This is only possible if $\{a_i\}$ and $\{b_i\}$ are orthonormal, or $\text{rank}(M)=1$

Given: $T = \sum a_i \otimes b_i \otimes c_i$

When can we recover the factors a_i , b_i and c_i uniquely?

Jennrich: If $\{a_i\}$ and $\{b_i\}$ are full rank and no pair in $\{c_i\}$ are scalar multiples of each other

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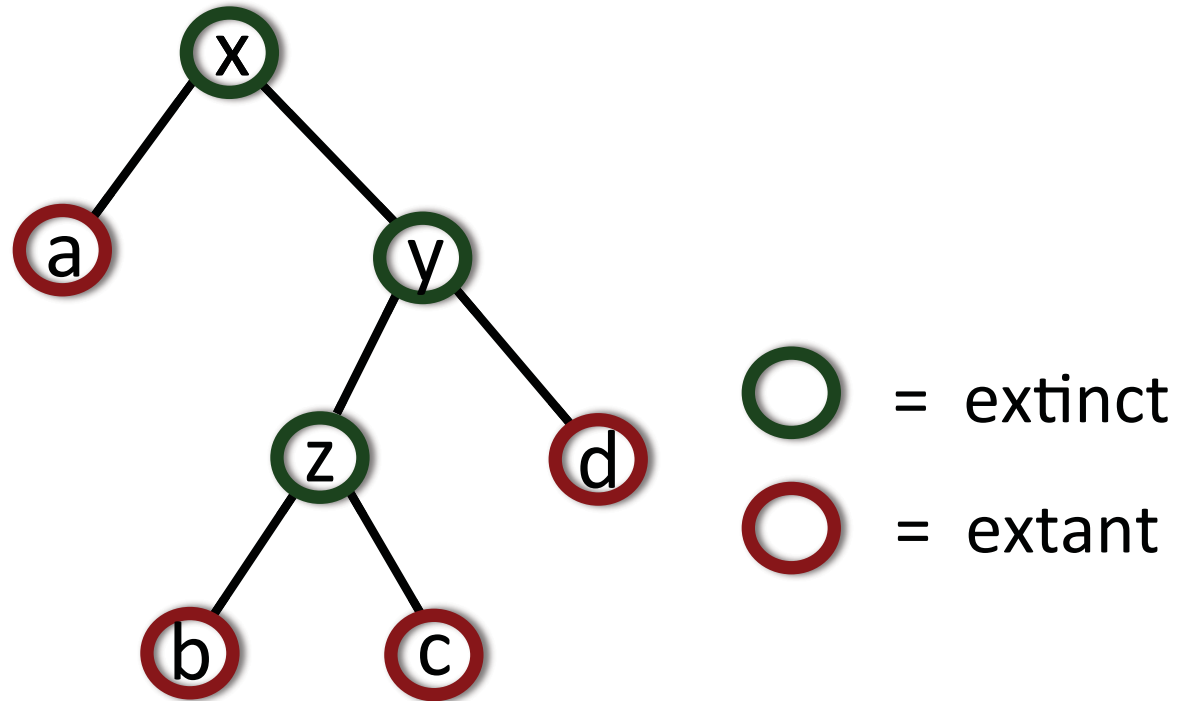
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Part III: Smoothed Analysis

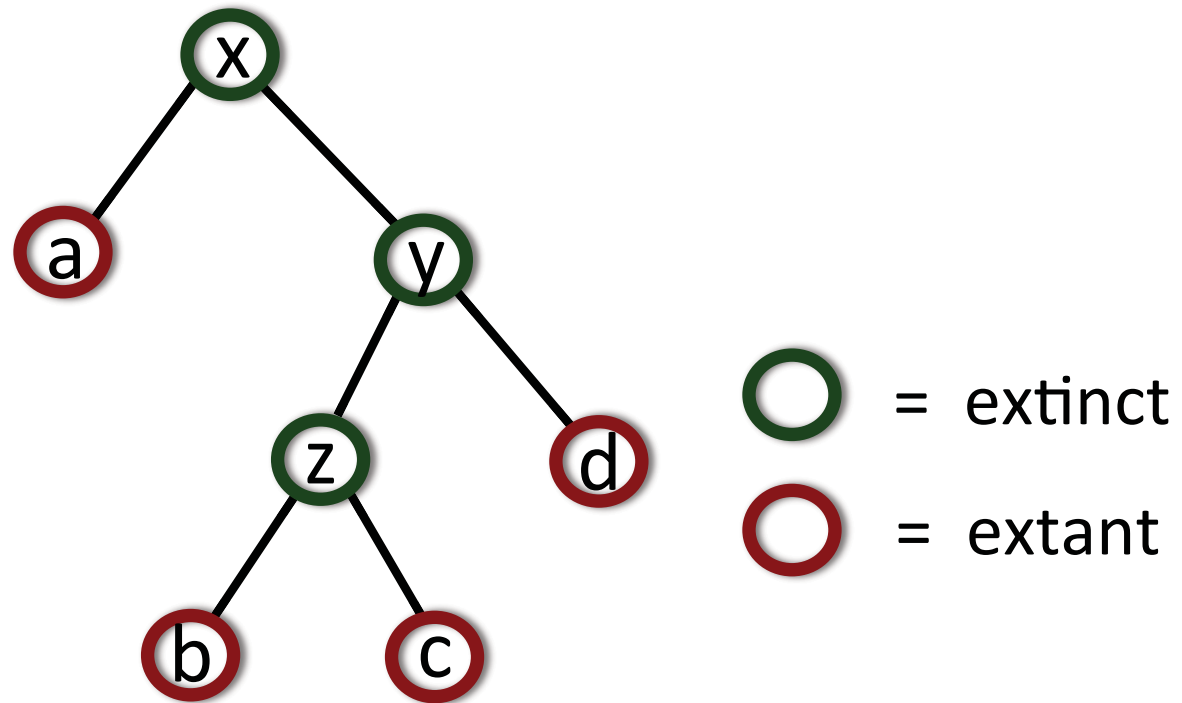
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PHYLOGENETIC RECONSTRUCTION

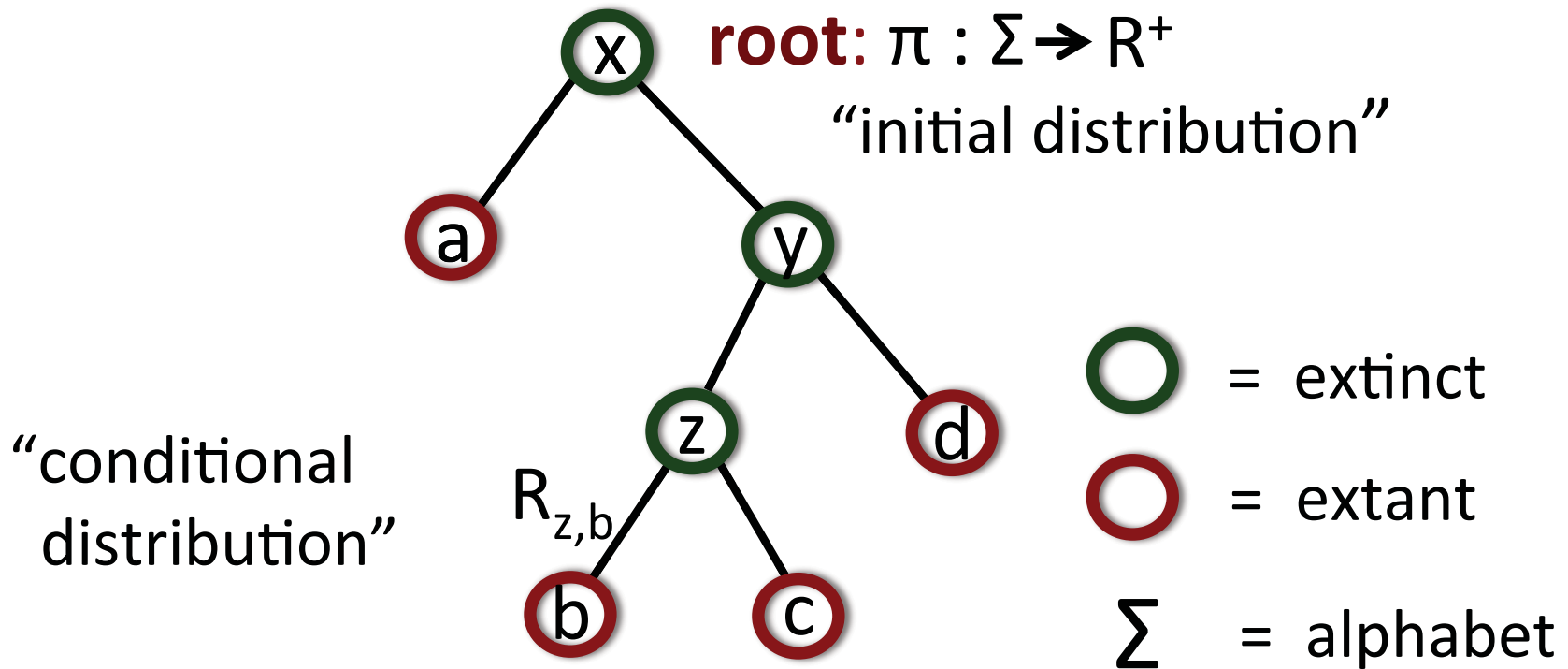


“Tree of Life”

PHYLOGENETIC RECONSTRUCTION



PHYLOGENETIC RECONSTRUCTION





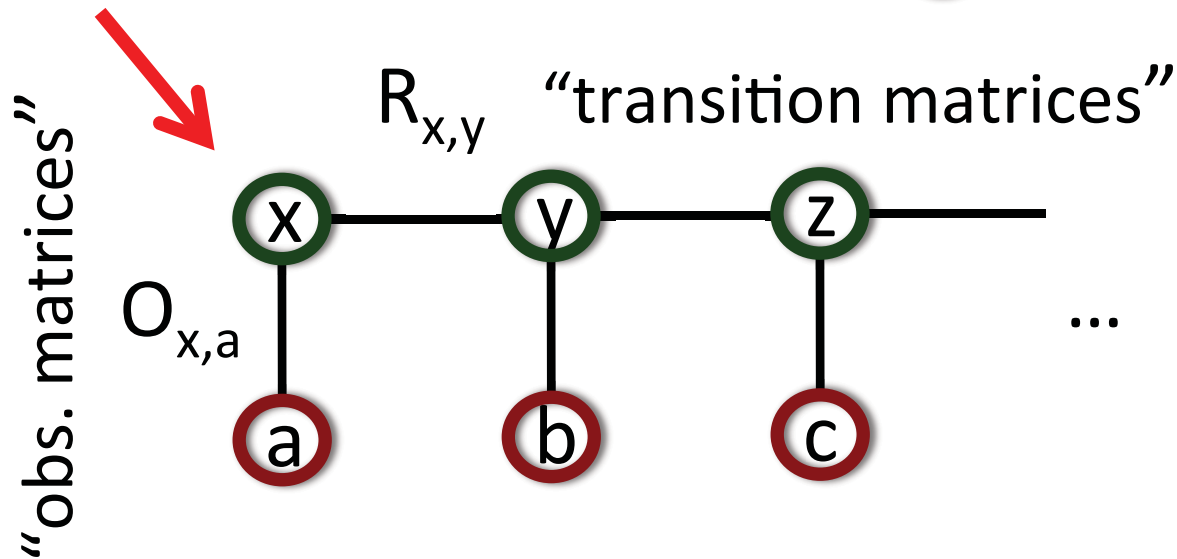
In each sample, we observe a symbol (Σ) at each extant (\bigcirc) node where we sample from π for the root, and propagate it using $R_{x,y}$, etc


HIDDEN MARKOV MODELS

$$\pi : \Sigma_s \rightarrow \mathbb{R}^+$$

“initial distribution”

 = hidden
 = observed



In each sample, we observe a symbol (Σ_o) at each obs. () node where we sample from π for the start, and propagate it using $R_{x,y}$, etc (Σ_s)

Question: Can we reconstruct just the topology from random samples?

Usually, we assume $T_{x,y}$, etc are full rank so that we can re-root the tree arbitrarily

[Steel, 1994]: The following is a distance function on the edges

$$d_{x,y} = -\ln |\det(P_{x,y})| + \frac{1}{2} \ln \prod_{\sigma \in \Sigma} \pi_{x,\sigma} - \frac{1}{2} \ln \prod_{\sigma \in \Sigma} \pi_{y,\sigma}$$

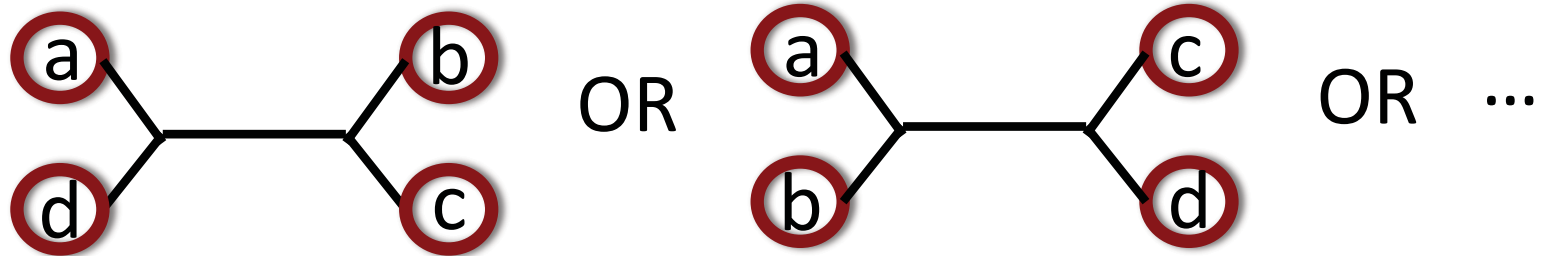
where $P_{x,y}$ is the joint distribution, and the distance between leaves is the sum of distances on the path in the tree

(It's not even obvious it's nonnegative!)

Question: Can we reconstruct just the topology from random samples?

Usually, we assume $T_{x,y}$, etc are full rank so that we can re-root the tree arbitrarily

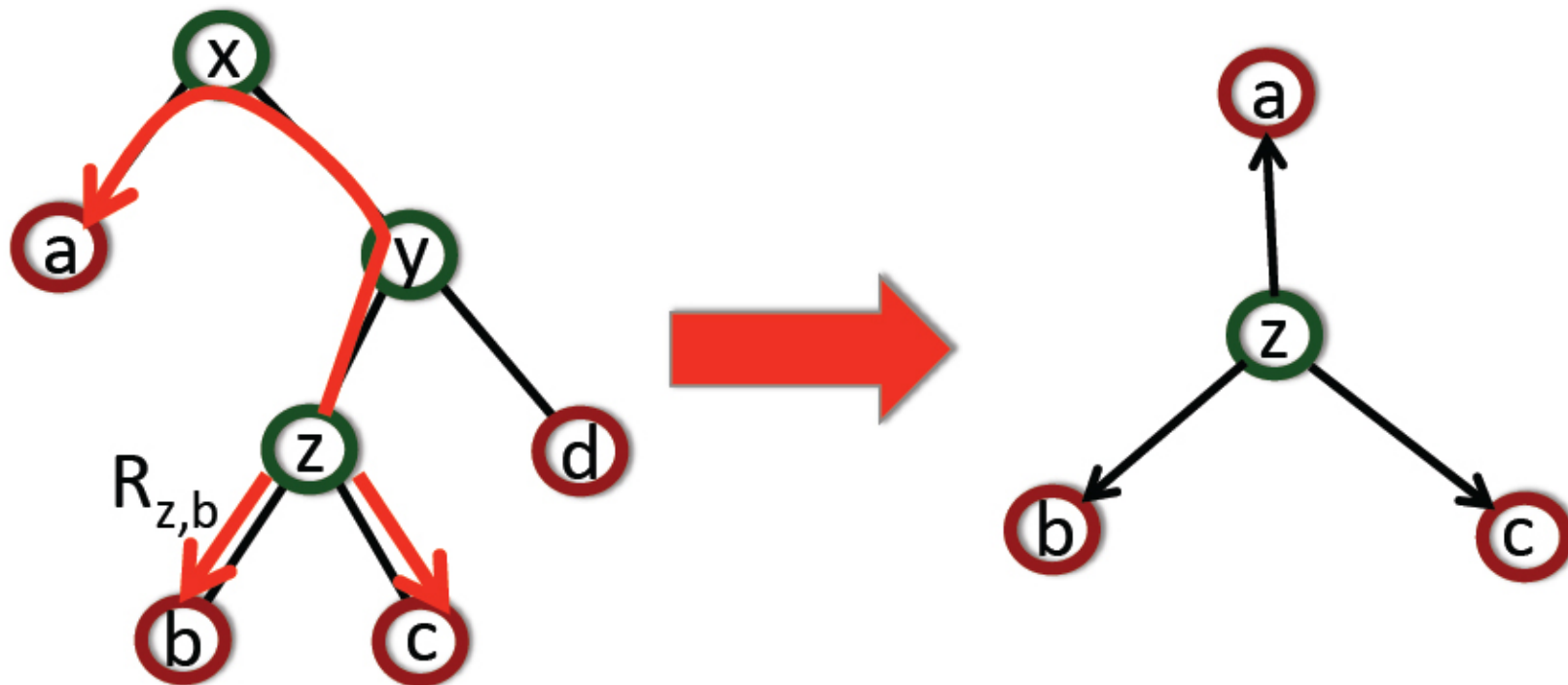
[Erdos, Steel, Szekely, Warnow, 1997]: Used Steel's distance function and quartet tests



to reconstruction the topology, from polynomially many samples

For many problems (e.g. HMMs) finding the transition matrices is the main issue...

[Chang, 1996]: The model is identifiable (if R's are full rank)



Joint distribution over (a, b, c):

$$\sum_{\sigma} \Pr[z = \sigma] \Pr[a | z = \sigma] \otimes \underbrace{\Pr[b | z = \sigma] \otimes \Pr[c | z = \sigma]}_{\text{columns of } R_{z,b}}$$

[Mossel, Roch, 2006]: There is an algorithm to PAC learn a phylogenetic tree or an HMM (if its transition/output matrices are full rank) from polynomially many samples

Question: Is the full-rank assumption necessary?

[Mossel, Roch, 2006]: It is as hard as noisy-parity to learn the parameters of a general HMM

Noisy-parity is an infamous problem in learning, where $O(n)$ samples suffice but the best algorithms run in time $2^{n/\log(n)}$

Due to **[Blum, Kalai, Wasserman, 2003]**

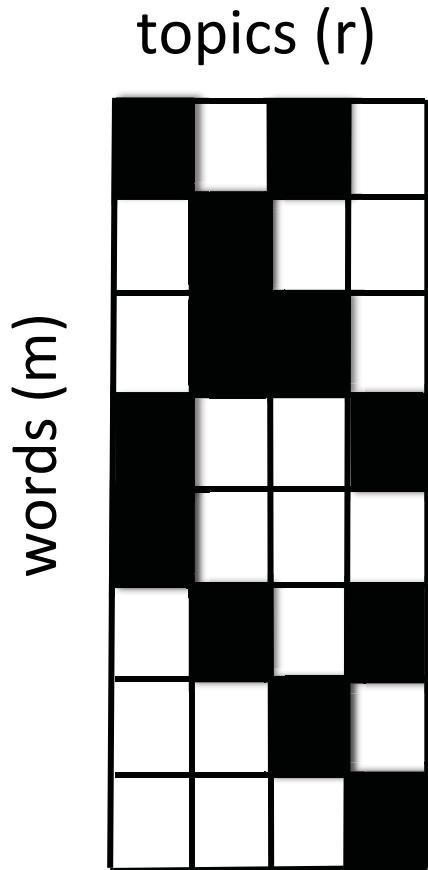
(It's now used as a hard problem to build cryptosystems!)

THE POWER OF CONDITIONAL INDEPENDENCE

[Phylogenetic Trees/HMMS]: (joint distribution on leaves a, b, c)

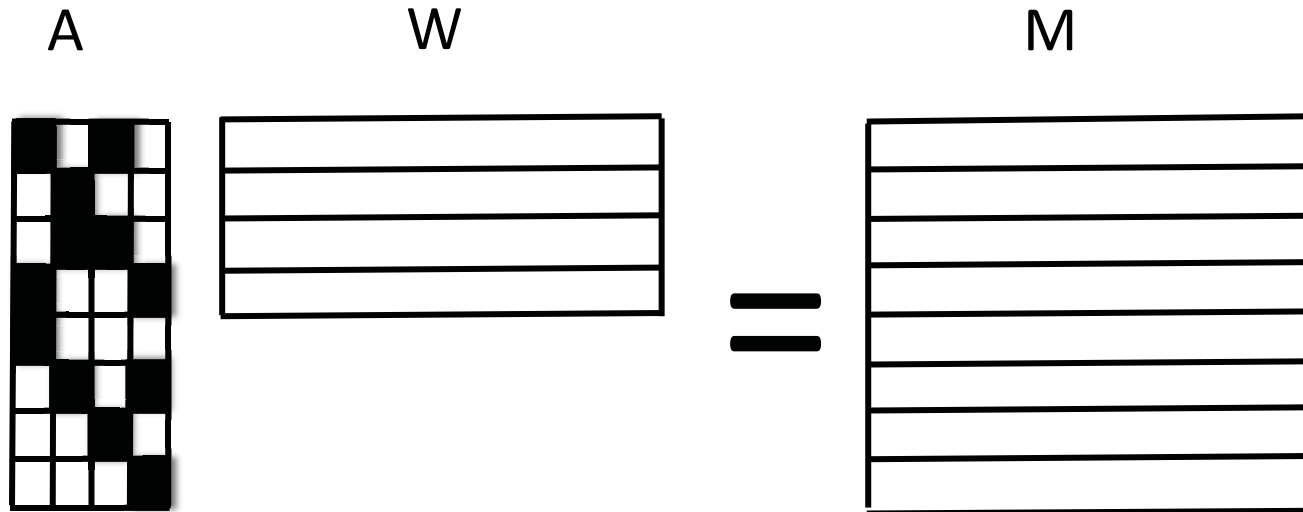
$$\sum_{\sigma} \Pr[z = \sigma] \Pr[a | z = \sigma] \otimes \Pr[b | z = \sigma] \otimes \Pr[c | z = \sigma]$$

PURE TOPIC MODELS

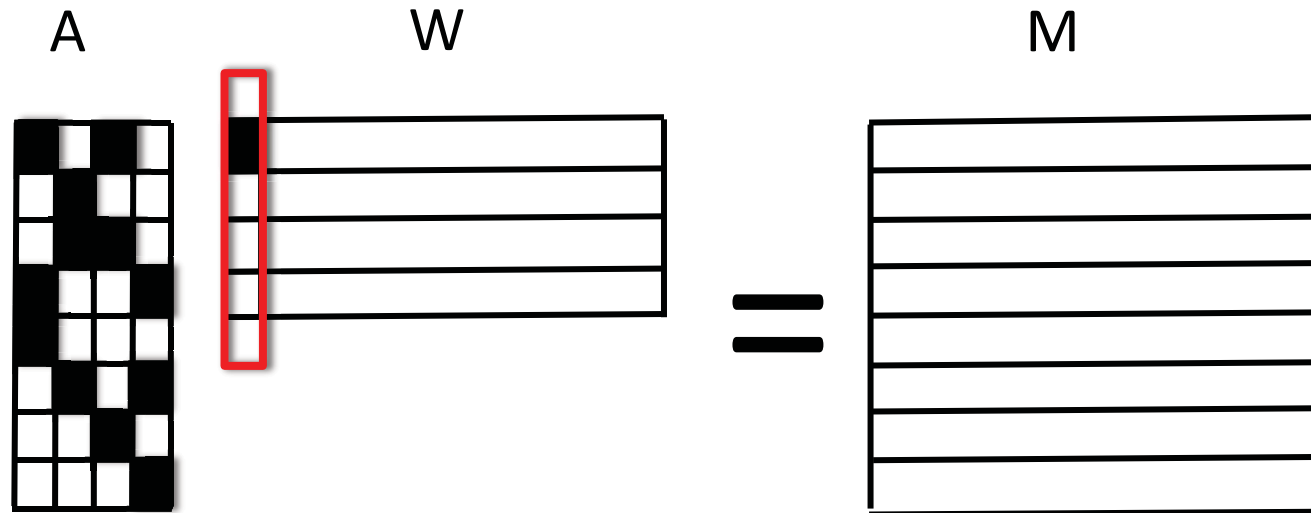


- Each topic is a distribution on words
- **Each document is about only one topic**
(stochastically generated)
- Each document, we sample L words from its distribution

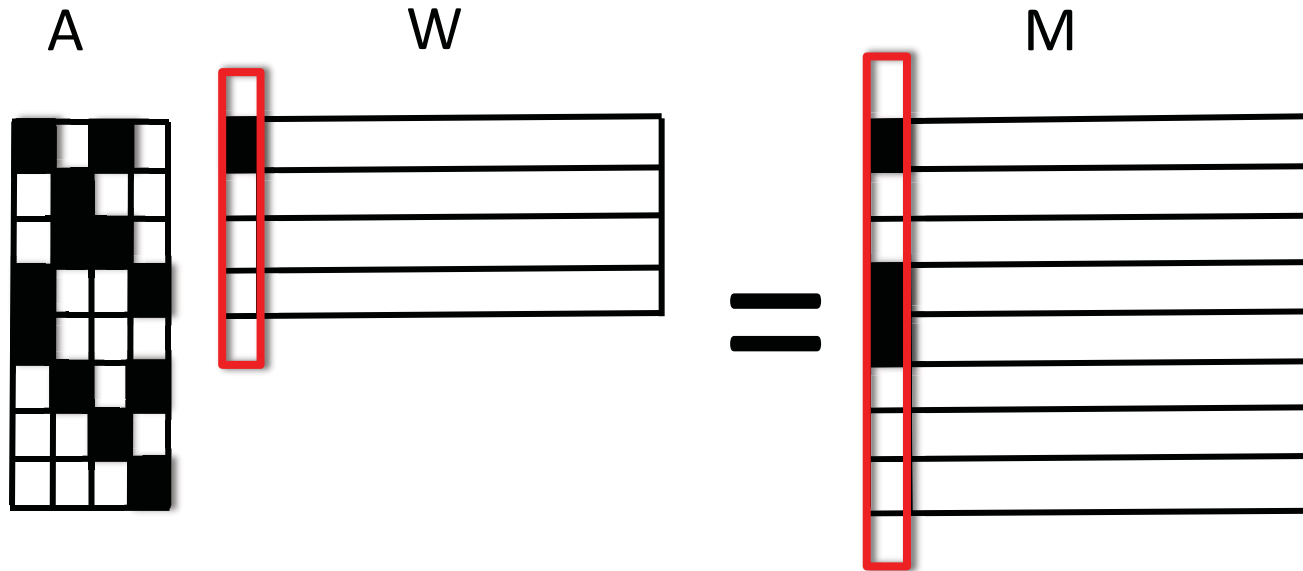
PURE TOPIC MODELS



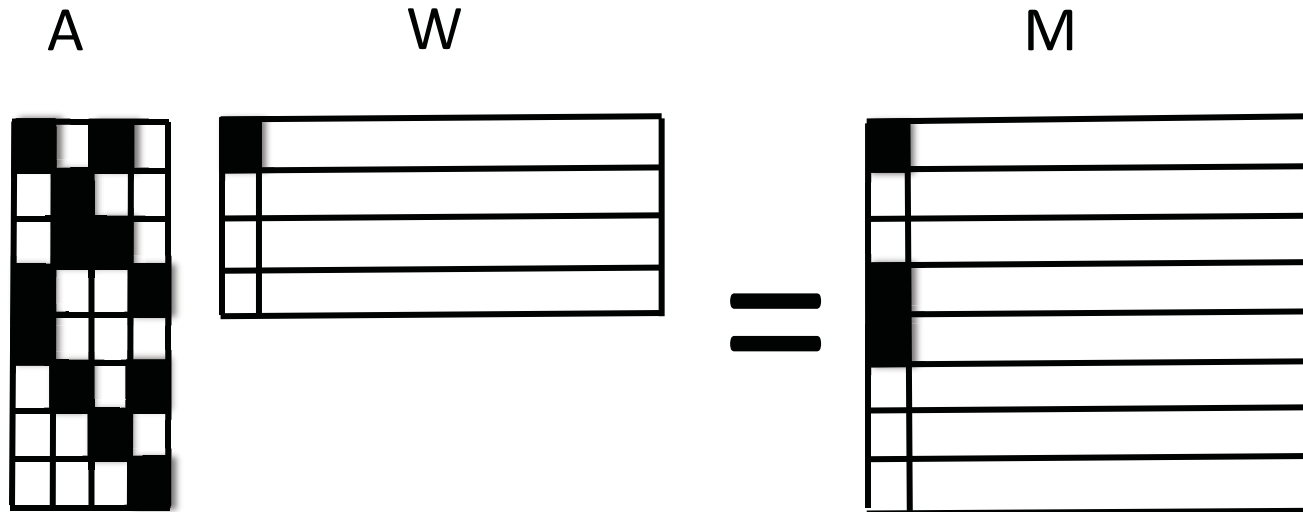
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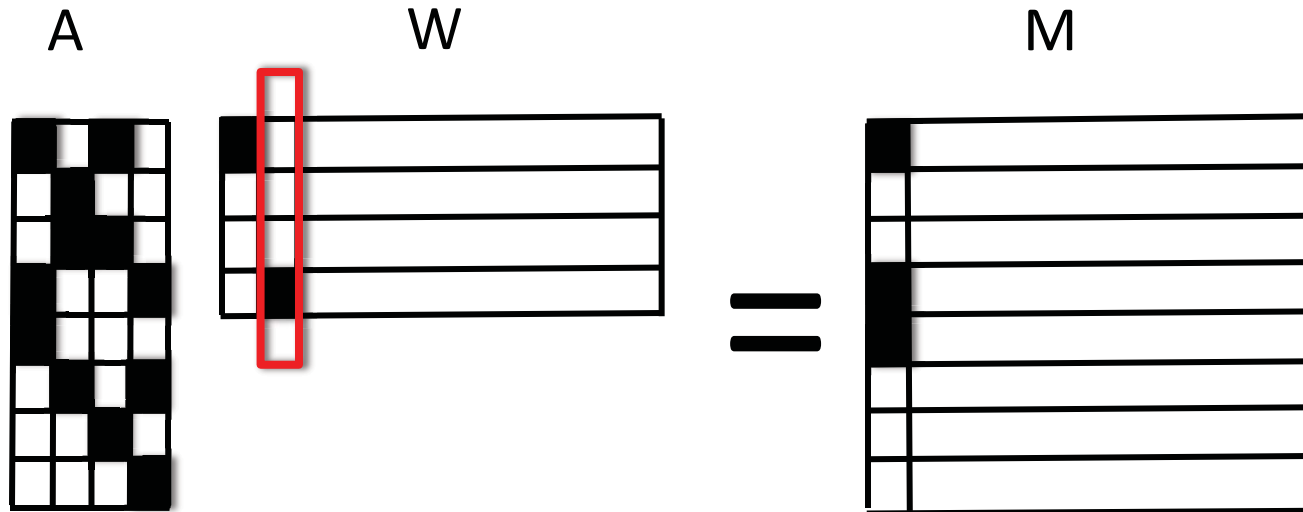
PURE TOPIC MODELS



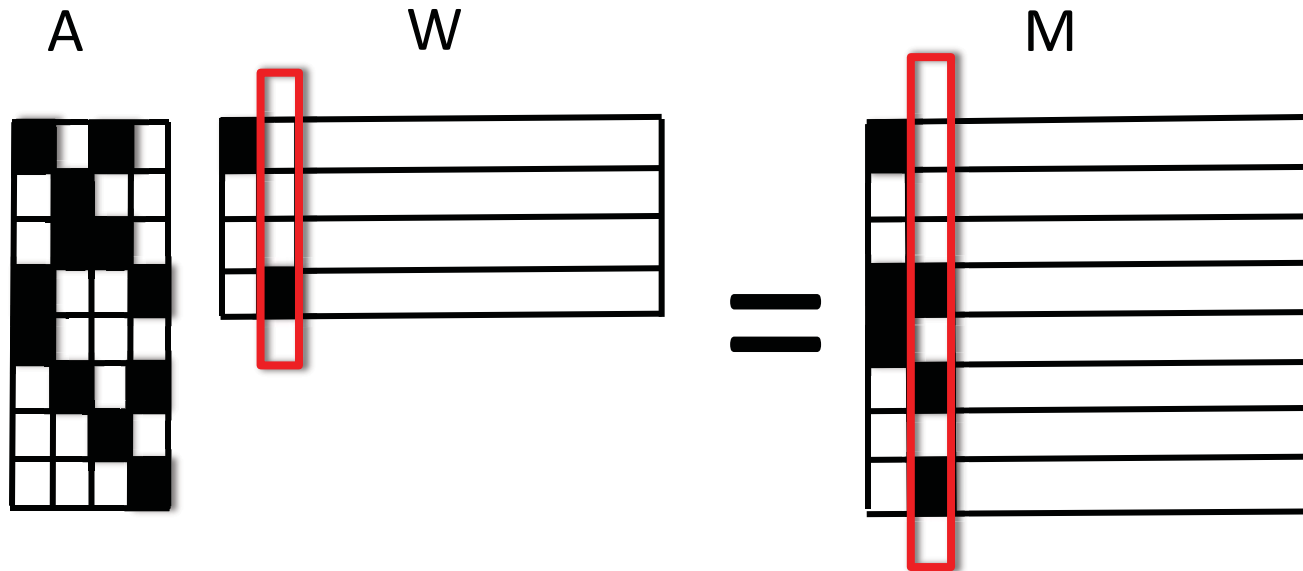
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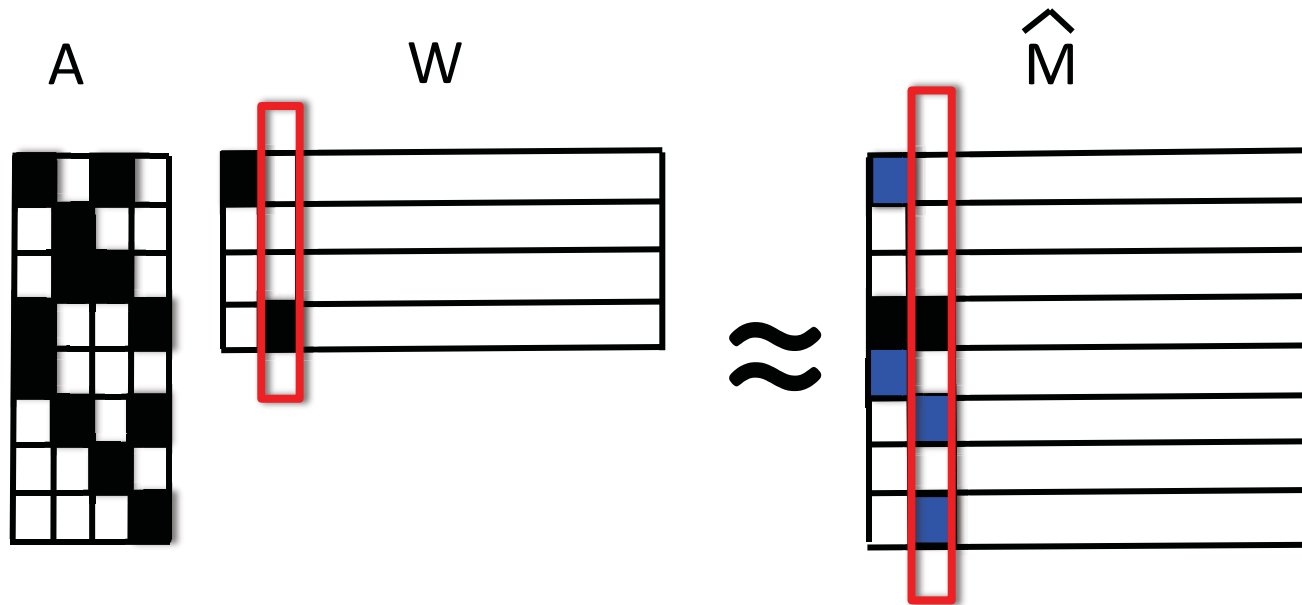
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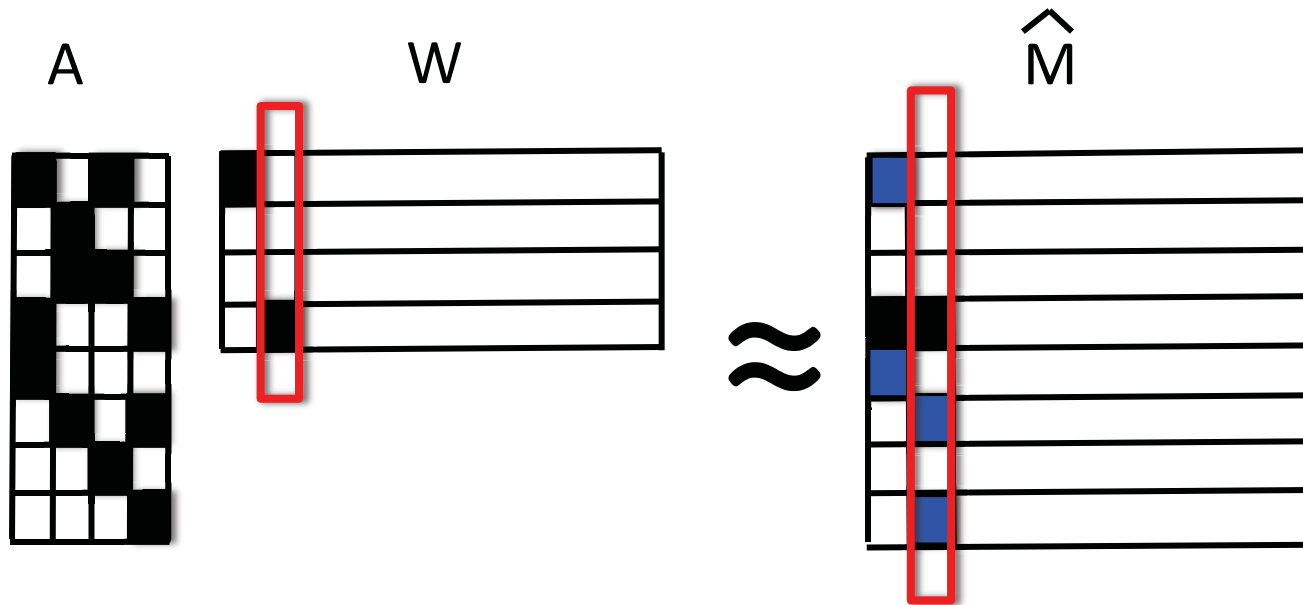
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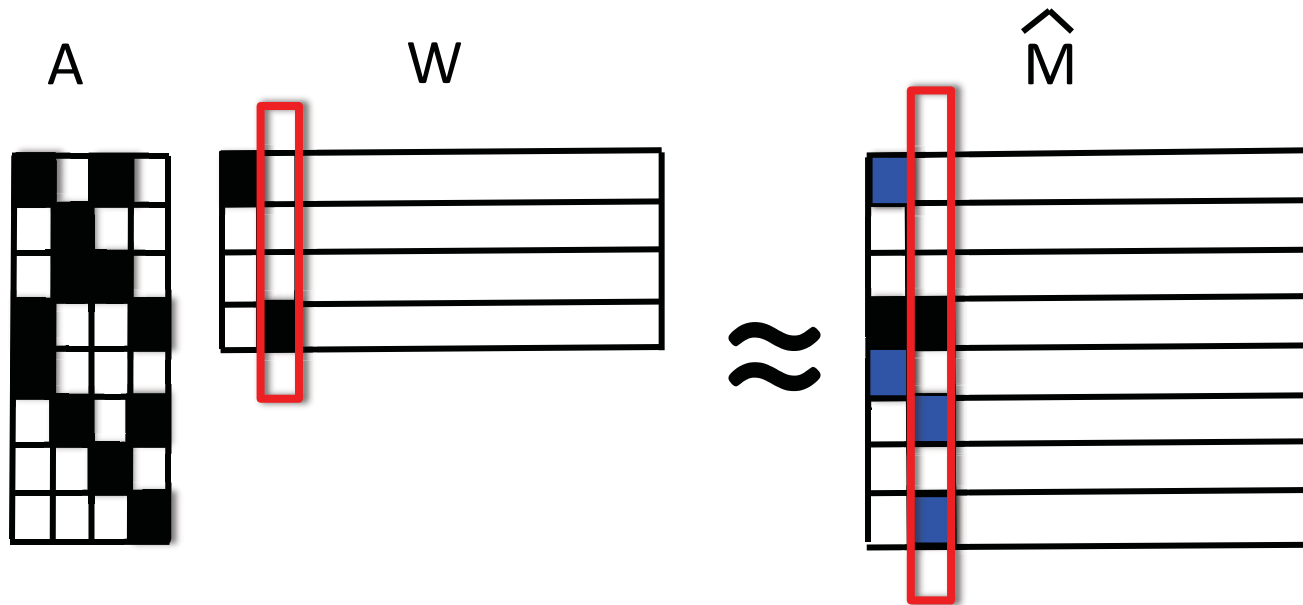


PURE TOPIC MODELS



[Anandkumar, Hsu, Kakade, 2012]: Algorithm for learning pure topic models from polynomially many samples (A is full rank)

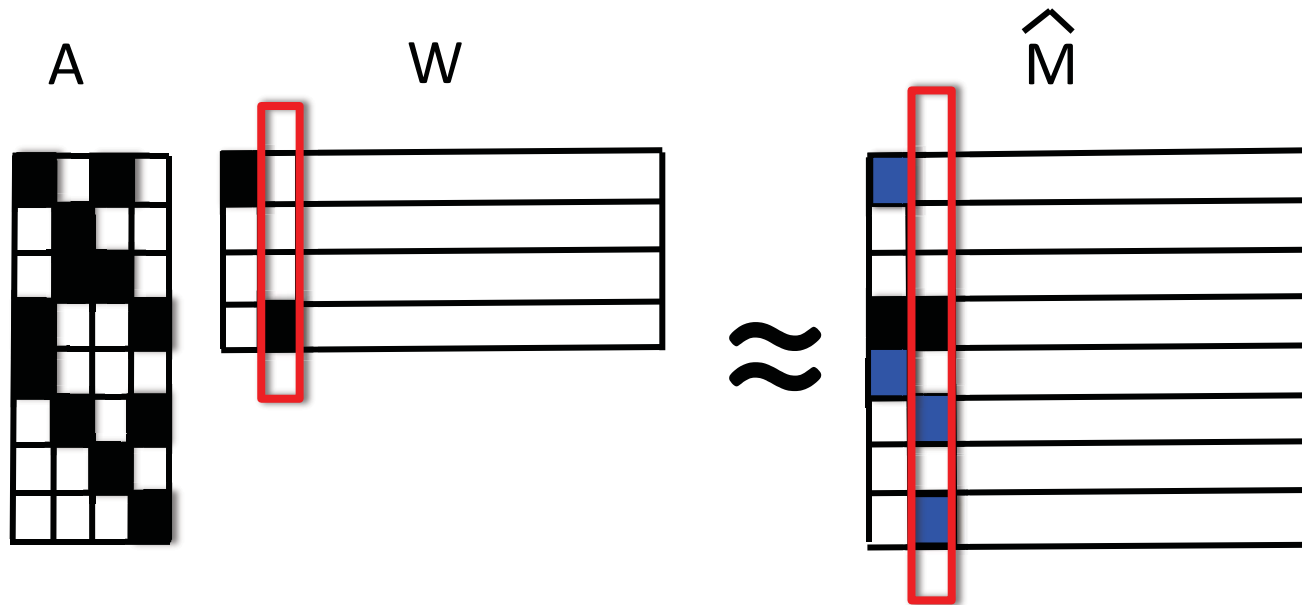
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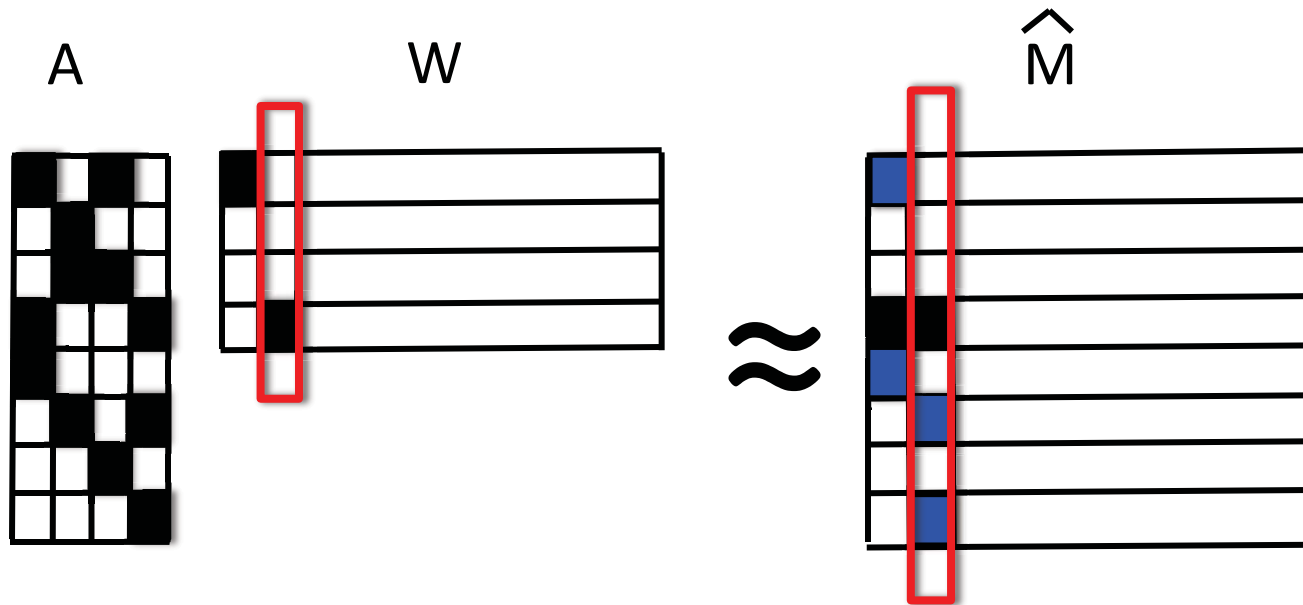
Question: Where can we find three conditionally independent random variables?

PURE TOPIC MODELS



[Anandkumar, Hsu, Kakade, 2012]: Algorithm for learning pure topic models from polynomially many samples (A is full rank)

PURE TOPIC MODELS



[Anandkumar, Hsu, Kakade, 2012]: Algorithm for learning pure topic models from polynomially many samples (A is full rank)

The first, second and third words are independent conditioned on the topic t (and are random samples from A_t)

THE POWER OF CONDITIONAL INDEPENDENCE

[Phylogenetic Trees/HMMS]: (joint distribution on leaves a, b, c)

$$\sum_{\sigma} \Pr[z = \sigma] \Pr[a | z = \sigma] \otimes \Pr[b | z = \sigma] \otimes \Pr[c | z = \sigma]$$

[Pure Topic Models/LDA]: (joint distribution on first three words)

$$\sum_j \Pr[\text{topic} = j] A_j \otimes A_j \otimes A_j$$

[Community Detection]: (counting stars)

$$\sum_j \Pr[C_x = j] (C_A \Pi)_j \otimes (C_B \Pi)_j \otimes (C_C \Pi)_j$$

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- Kruskal Rank and the Khatri-Rao Product

So far, Jennrich's algorithm has been the key but it has a crucial limitation. Let

$$T = \sum_{i=1}^R a_i \otimes a_i \otimes a_i$$

where $\{a_i\}$ are n -dimensional vectors

Question: What if R is much larger than n ?

This is called the **overcomplete** case — e.g. the number of factors is much larger than the number of observations...


In such cases, why stop at third-order tensors?

Consider a **sixth**-order tensor T:

$$T = \sum_{i=1}^R a_i \otimes a_i \otimes a_i \otimes a_i \otimes a_i \otimes a_i$$

Question: Can we find its factors, even if R is much larger than n?

Let's flatten it by rearranging its entries into a **third**-order tensor:

$$\text{flat}(T) = \sum_{i=1}^R b_i \otimes b_i \otimes b_i \quad (\text{where } b_i = a_i \otimes_{KR} a_i)$$


n^2 -dimensional vector whose $(j,k)^{\text{th}}$ entry is the product of the j^{th} and k^{th} entries of a_i — **Khatri-Rao product**

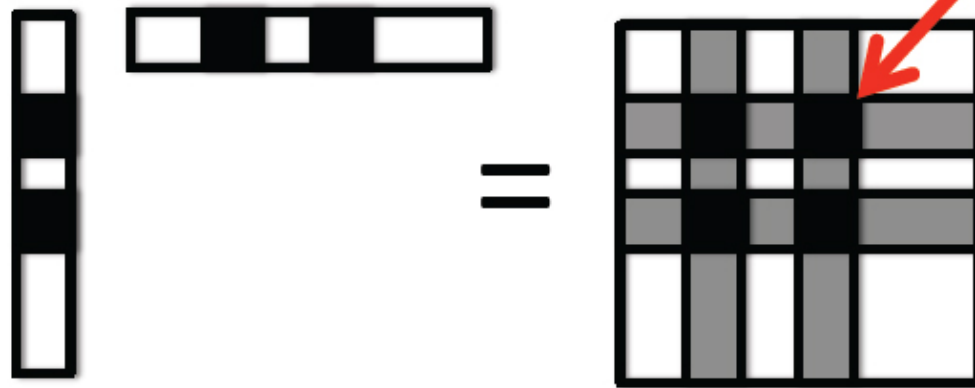
Question: Can we apply Jennrich's Algorithm to $\text{flat}(T)$?

When are the new factors $b_i = a_i \otimes_{\mathbb{K}\mathbb{R}} a_i$ linearly independent?

Example #1:

Let $\{a_i\}$ be all $\binom{n}{2}$ vectors with exactly two ones

Then $\{b_i\}$ are vectorizations of:



and are linearly independent

Question: Can we apply Jennrich's Algorithm to $\text{flat}(T)$?

When are the new factors $b_i = a_i \otimes_{KR} a_i$ linearly independent?

Example #2:

Let $\{a_{1..n}\}$ and $\{a_{n+1..2n}\}$ be two random orthonormal bases

Then there is a linear dependence with $2n$ terms:

$$\sum_{i=1}^n a_i \otimes_{KR} a_i - \sum_{i=n+1}^{2n} a_i \otimes_{KR} a_i = 0$$

(as matrices, both sum to the identity)

THE KRUSKAL RANK

Definition: The **Kruskal rank** (k-rank) of $\{b_i\}$ is the largest k s.t. every set of k vectors is linearly independent

$$b_i = a_i \underset{KR}{\otimes} a_i \quad k\text{-rank}(\{a_i\}) = n$$

Example #1: $k\text{-rank}(\{b_i\}) = R = \binom{n}{2}$

Example #2: $k\text{-rank}(\{b_i\}) = 2n-1$

The Kruskal rank always **adds** under the Khatri-Rao product, but sometimes it **multiplies** and that can allow us to handle $R \gg n$

[Allman, Matias, Rhodes, 2009]: Almost surely, the Kruskal rank multiplies under the Khatri-Rao product

Proof: The set of $\{a_i\}$ where

$$b_i = a_i \otimes_{KR} a_i \quad \text{and} \quad \det(\{b_i\}) = 0$$

is measure zero ■

But this yields a very weak bound on the **condition number** of $\{b_i\}$...

... which is what we need to apply it to learning/statistics, where we have an estimate to T

[Allman, Matias, Rhodes, 2009]: Almost surely, the Kruskal rank multiplies under the Khatri-Rao product

Definition: The **robust Kruskal rank** ($k\text{-rank}_\gamma$) of $\{b_i\}$ is the largest k s.t. every set of k vector has condition number at most $O(\gamma)$

[Bhaskara, Charikar, Vijayaraghavan, 2013]: The robust Kruskal rank always under the Khatri-Rao product

[Bhaskara, Charikar, Moitra, Vijayaraghavan, 2014]: Suppose the vectors $\{a_i\}$ are ε -perturbed. Then

$$k\text{-rank}_\gamma(\{b_i\}) = R$$

for $R = n^2/2$ and $\gamma = \text{poly}(1/n, \varepsilon)$ with **exponentially** small failure probability (δ)

[Bhaskara, Charikar, Moitra, Vijayaraghavan, 2014]: Suppose the vectors $\{a_i\}$ are ε -perturbed. Then

$$k\text{-rank}_\gamma(\{b_i\}) = R$$

for $R = n^2/2$ and $\gamma = \text{poly}(1/n, \varepsilon)$ with **exponentially** small failure probability (δ)

Hence we can apply Jennrich's Algorithm to $\text{flat}(T)$ with $R \gg n$

Note: These bounds are easy to prove with inverse **polynomial** failure probability, but then γ depends δ

This can be extended to any constant order Khatri-Rao product

[Bhaskara, Charikar, Moitra, Vijayaraghavan, 2014]: Suppose the vectors $\{a_i\}$ are ε -perturbed. Then

$$k\text{-rank}_\gamma(\{b_i\}) = R$$

for $R = n^2/2$ and $\gamma = \text{poly}(1/n, \varepsilon)$ with **exponentially** small failure probability (δ)

Hence we can apply Jennrich's Algorithm to $\text{flat}(T)$ with $R \gg n$

Sample application: Algorithm for learning mixtures of $n^{O(1)}$ spherical Gaussians in R^n , if their means are ε -perturbed

This was also obtained independently by **[Anderson, Belkin, Goyal, Rademacher, Voss, 2014]**

Any Questions?

Summary:

- Tensor decompositions are **unique** under much more general conditions, compared to matrix decompositions
- Jennrich's Algorithm (rediscovered many times!), and its many applications in learning/statistics
- Introduced **new models** to study overcomplete problems ($R \gg n$)
- Are there algorithms for order- k tensors that work with $R = n^{0.51 k}$?

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